

Long Information Design

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Work on *zero-sum dynamic splitting games*.

Two information designers have access to independent sources of information and want to influence a decision-maker.

Outline:

1. Introduction
2. Model
3. Games with deadline
4. Games without deadline

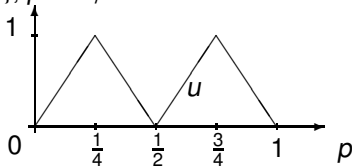
Introduction

Bayesian Persuasion

Some k in K is the true state. Public a priori p in $\Delta(K)$

A player chooses M and $x : K \rightarrow \Delta(M)$. Then k is chosen ac. to p , a message m is chosen ac. to $x(k)$ and publicly sent to decision-makers. Payoff to the player: $u(p(m))$, where $p(m)$ is the posterior on K given m , and $u : \Delta(K) \rightarrow \mathbb{R}$ is continuous.

Example: $K = \{0, 1\}$, $p = 1/2$.

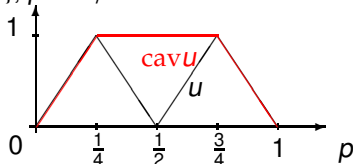


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The value is $\text{cav}u(p) = \inf\{f(p), f \text{ concave}, f \geq u\}$

$$= \max \left\{ \sum_{m \in K} \lambda_m u(p_m), \forall m \lambda_m \geq 0, p_m \in \Delta(K), \sum_m \lambda_m = 1, \sum_m \lambda_m p_m = p \right\}.$$

(Aumann Maschler 1966. Kamenica Gentzkow 2011)

Sometimes not all splittings are available

Example: $K = \{0, 1\}$. The player (pharmaceutical firm) can run one or several experiments, each experiment returns the true state with proba $2/3$. Assume $p_0 = 1/2$ and $u(p) = \mathbf{1}_{p \geq 4/5}$ for all p .

An experiment from p_0 induces the distribution $\frac{1}{2}\delta_{2/3} + \frac{1}{2}\delta_{1/3}$ on $\Delta(K)$. From p , it induces

$$\varphi(p) = \frac{(1+p)}{3}\delta_{2p/(1+p)} + \frac{(2-p)}{3}\delta_{p/(2-p)}.$$

Optimum: run several experiments, stop if the current belief is $\geq 4/5$.

Say that f is S -concave if: $\forall p, f(p) \geq \frac{1}{3}(1+p)f(\frac{2p}{1+p}) + \frac{1}{3}(2-p)f(\frac{p}{2-p})$.

Theorem : the value is $v(p) = \inf\{f(p), f \text{ } S\text{-concave}, f \geq u\}$

(*Fundamental theorem of Gambling*, Dubbins and Savage 1965)

Here $v(p_0) = 5/8$. v is not continuous on $[0, 1]$.

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We now study **competition between 2 information designers**, who control independent information.

Example: Competing firms release information about a new product to influence buyers' valuations.

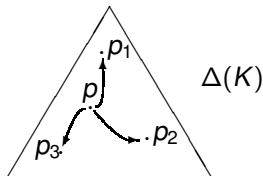
We use tools from the literature on dynamic games: Mertens-Zamir (1971), Heuer (1992), Renault-Venel (2017), Oliu-Barton (2018), Laraki-Renault (2020). Name reminiscent of “Long Cheap Talk” (Aumann Hart 2003).

Here : terminal payoffs and constraints on splittings.

Notations:

For $p \in \Delta(K)$, the set of all possible splittings of p is

$$SPL(p) = \left\{ s \in \Delta(\Delta(K)) : \int_{p' \in \Delta(K)} p' ds(p') = p \right\}.$$



Given $u : \Delta(K) \times \Delta(L) \rightarrow \mathbb{R}$ measurable bounded, we extend u to distributions by:

$$u(s, t) = \int_{(p, q) \in \Delta(K) \times \Delta(L)} u(p, q) ds(p) dt(q).$$

Model

- ▶ Finite set of unknown states $K \times L$, compact subsets P of $\Delta(K)$ and Q of $\Delta(L)$. Continuous payoff function $u : P \times Q \rightarrow \mathbb{R}$.
- ▶ Player 1 designs information about K , controls belief p in P . Available splittings at p : $S(p) \subset SPL(p) \cap \Delta(P)$, with $\delta_p \in S(p)$.

Player 2 designs information about L , controls belief q in Q .

Available splittings at q : $T(q) \subset SPL(q) \cap \Delta(Q)$, with $\delta_q \in T(q)$.

S and T : continuous with convex compact values.

- ▶ At the end of each period, past splittings and posteriors are publicly observed.
- ▶ Finally, an agent takes an optimal decision as a function of his final posterior beliefs (p, q) , inducing an expected payoff $u(p, q)$ for player 1 and $-u(p, q)$ for player 2.

Timing: we consider several game forms.

- ▶ T -period simultaneous game G_T and infinite horizon game G_∞ :
In each period t , given the current beliefs (p_{t-1}, q_{t-1}) , simultaneously player 1 chooses a splitting in $S(p_{t-1})$ and player 2 chooses a splitting in $T(q_{t-1})$.
- ▶ T -period sequential games H_T and infinite horizon game H_∞ :
Player 1 splits in odd periods, player 2 in even periods.

Examples of admissible splittings:

- ▶ Standard case: player 1 can choose any experiment $x : K \rightarrow \Delta(M)$. Then $S(p) = SPL(p)$ for all p .
- ▶ More generally: player 1 has access to all experiments in a given compact subset X of $\{x : K \rightarrow \Delta(M)\}$.
- ▶ Player 1 can choose experiments with posteriors in a finite set A .

S-concavification and T-convexification

Let $f : P \times Q \rightarrow \mathbb{R}$.

Say that f is S-concave if for all $p, q, s \in S(p)$: $f(s, q) \leq f(p, q)$.

Say that f is T-convex if for all $p, q, t \in T(q)$: $f(p, t) \geq f(p, q)$.

$$\text{cav}f(p, q) := \inf\{g(p, q), g(\cdot, q) \text{ concave}, g \geq f\},$$

$$\text{vex}f(p, q) := \sup\{g(p, q), g(p, \cdot) \text{ convex}, g \leq f\}.$$

Assumption 1 on available splittings:

Any distribution which can be reached in 2 stages can be reached in a single stage: $S^2 = S$ and $T^2 = T$.

Then for f continuous:

$$\begin{aligned} \text{cav}f(p, q) &= \max\{f(s, q), s \in S(p)\}, \\ \text{vex}f(p, q) &= \min\{f(p, t), t \in T(q)\}. \end{aligned}$$

Games with deadline

Proposition 1: Under A1,

Finite Sequential Games:

The game H_1 has value $\text{cav}u(p_0, q_0) = \max_{s \in SP(p_0)} u(s, q_0)$,

The game H_2 has value $\text{cav} \text{vex}u(p_0, q_0)$ (same for all games H_{2T}),

The game H_3 has value $\text{vex} \text{cav}u(p_0, q_0)$ (same for all games H_{2T+1}),

Finite Simultaneous Games:

G_1 has value

$$\text{OSV}(p_0, q_0) = \max_{s \in S(p_0)} \min_{t \in S(q_0)} u(s, t) = \min_{t \in S(q_0)} \max_{s \in S(p_0)} u(s, t)$$

One-shot splitting value (Laraki, Sorin 2001 for the unconstrained case), may not exist if u is discontinuous. Same value for all games G_T .

Remark: The values (except for H_1) are all concave-convex. *It is always safe to remain silent if you can transmit information later.*

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Games without deadline

We use $\|\cdot\|_1$ on P and $\|\cdot\|_\infty$ on $C(P, \mathbb{R})$.

Let us introduce the following set (from R-Venel 2017) :

$$D = \{f : P \rightarrow [-1, 1], \forall p, p' \in P, \forall a, b \geq 0, |af(p) - bf(p')| \leq \|ap - bp'\|\}.$$

If $f \in D$, f is 1-Lispchitz.

D is a convex lattice, containing all constants in $[-1, 1]$ and linear forms on \mathbb{R}^K with norm ≤ 1 .

Lemma The set $L :=_{\text{def}} \{\alpha f, \alpha \geq 0, f \in D\}$ is dense in $C(P, \mathbb{R})$.

Notice that L contains all non revealing functions à la Aumann Maschler.

Assumption A2 on available splittings: $\forall f \in D, \text{cav}f \in D$
(and similarly for player 2)

A2 is equivalent to: $\forall p, p' \in P, \forall s \in S(p), \forall a, b \geq 0,$

$$\exists s' \in S(p') \text{ s.t. } \forall f \in D, |af(s) - bf(s')| \leq \|ap - bp'\|,$$

and similarly for T .

(It implies that $S : P \rightrightarrows \Delta(P)$ is non expansive for the distance $d^*(u, v) = \sup_{f \in D} |u(f) - v(f)|$. Close to a condition in R-Venel 17).

Illustration: Fix a single experiment $x : K \rightarrow \Delta(M)$. $p \in \Delta(K)$ induces the distribution:

$$\varphi(p) = \sum_{m \in M} x(p)(m) \delta_{l(p,m)},$$

with $x(p)(m) = \sum_k p^k x^k(m)$ probability of m , and $l(p, m) = \left(\frac{p^k x^k(m)}{x(p)(m)} \right)_k$ conditionnal probability on K given m .

For f in D , p, p' in P , $a, b \geq 0$,

$$\begin{aligned} a f(\varphi(p)) - b f(\varphi(p')) &= \sum_{m \in M} a x(p)(m) f(l(p, m)) - b x(p')(m) f(l(p', m)), \\ &\leq \sum_{m \in M} \|a (p^k x^k(m))_k - b (p'^k x^k(m))_k\|, \\ &= \sum_{m \in M} \sum_{k \in K} |a p^k x^k(m) - b p'^k x^k(m)| = \|ap - bp'\|. \end{aligned}$$

and A2 is satisfied.

(can be extended to M compact)

(in general we don't have here: f 1-Lipschitz $\implies \text{cav} f$ is 1-Lipschitz)

Example: Assume player 1 has access to all experiments in a given compact subset X of $\{x : K \rightarrow \Delta(M)\}$ including the non revealing experiment.

For any “independent” strategy $\sigma = (\sigma_t)_{t \geq 1}$ with σ_t measurable from $(X \times M)^{t-1}$ to $\Delta(X)$, and $T \geq 1$, denote by $\mu_T(p, \sigma)$ the induced law of the posteriors if σ is played up to T .

Define $S_T(p) = \{\mu_T(p, \sigma), \sigma \text{ varies}\}$ and finally $S(p)$ as the closure of $\cup_{t \geq 1} S_T(p)$.

S satisfies the assumptions A1 and A2.

Example: Assume P, Q finite. A1 and A2 are satisfied (up to a small modification in the definition of D).

Theorem: Under A1 and A2,

1) There exists a unique continuous function $v : P \times Q \rightarrow \mathbb{R}$ which is S -concave, T -convex, and such that for all $(p, q) \in P \times Q$:

Q1: if $v(p, q) > u(p, q)$, there exists $s \in S(p)$ s.t. $v(p, q) = v(s, q)$ and $v(p', q) \leq u(p', q)$ for all $p' \in \text{supp}(s)$,

Q2: if $v(p, q) < u(p, q)$, there exists $t \in T(q)$ s.t. $v(p, q) = v(p, t)$ and $v(p, q') \geq u(p, q')$ for all $q' \in \text{supp}(t)$.

v is characterized by $v = \text{cav}_S \min(u, v) = \text{vex}_T \max(u, v)$.

We write $v = \text{MZ}(u)$, the Mertens-Zamir value induced by u .

2) Both infinite horizon games G_∞ and H_∞ have pure stationary optimal strategies, and value: $\text{MZ}(u)(p_0, q_0)$.

Remark : for G_∞, H_∞ , same as if the game stops after 2 silent moves.

Elements of Proof:

1) **Uniqueness of $MZ(u)$** : as in Laraki-R 2020.

Existence of $MZ(u)$: Let D_2 be the set of f in $C(P \times Q)$ such that
 $\forall p, p' \in P, \forall q, q' \in Q, \forall a, b \geq 0,$
 $|af(p, q) - bf(p', q)| \leq \|ap - bp'\|$ and $|af(p, q) - bf(p, q')| \leq \|aq - bq'\|.$

a) First assume that $u \in D_2$.

a1) Show that for each $\lambda \in (0, 1]$, the value v_λ of the λ discounted game is in D_2 .

a2) Consider a limit point $v = \lim_n v_{\lambda_n}$ and show that v is S -concave, T convex and satisfies Q1 and Q2. So $v = MZ(u)$. (close to Laraki-R 2020)

b) Conclude by density of $\{\alpha f, \alpha \geq 0, f \in D_2\}$ and the fact that $(u \mapsto MZ(u))$ is non expansive on D_2 .

2) In the game G_∞ or H_∞ , consider the strategy of player 1:
 At (p, q) , play $s \in S(p)$ given by Q1 (do nothing if $u(p, q) \geq v(p, q)$).
 Consider any strategy of player 2.

$(p_n, q_n)_n$ is a bounded martingale, hence CV a.s. to some (p_∞, q_∞) .
 We want to show that $\mathbb{E}(u(p_\infty, q_\infty)) \geq v(p_0, q_0)$.

For $n \geq 0$, $\mathbb{E}(v(p_{n+1}, q_n) | (p_n, q_n)) = v(p_n, q_n)$ by (Q1),

and $\mathbb{E}(v(p_{n+1}, q_{n+1}) | (p_{n+1}, q_n)) \geq v(p_{n+1}, q_n)$ since v is T -convex.

So

$$\mathbb{E}(v(p_{n+1}, q_{n+1})) \geq \mathbb{E}(v(p_{n+1}, q_n)) = \mathbb{E}(v(p_n, q_n)) \geq v(p_0, q_0).$$

By (Q1) again, $(p_{n+1}, q_n) \in \{u \geq v\}$ for each n , so passing to the limit we obtain $u(p_\infty, q_\infty) \geq v(p_\infty, q_\infty)$ a.s.

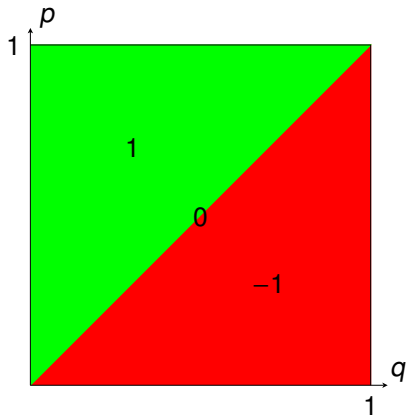
$$\begin{aligned} \mathbb{E}(u(p_\infty, q_\infty)) &\geq \mathbb{E}(v(p_\infty, q_\infty)) = \mathbb{E}(\lim_n (v(p_n, q_n))) \\ &= \lim_n \mathbb{E}(v(p_n, q_n)) \geq v(p_0, q_0). \end{aligned}$$

Remark: Generalized to a specific case of discontinuous u .

Illustration 1: being perceived as better

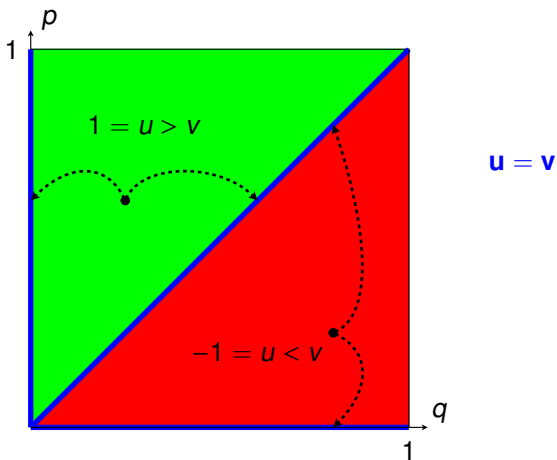
$K = L = \{0, 1\}$, all splittings are available. $p_0 = q_0 = 1/2$.

Discontinuous u function:



One-shot splitting value (Boleslavsky and Cotton, 2015) :

$OSV(p_0, q_0) = 0$ and it is optimal to split uniformly on $[0, 1]$.



$$v(p, q) = \frac{p - q}{\max(p, q)}$$

At equilibrium, the martingale moves at most once.

Illustration 2: Splitting on a finite grid

Fix P and Q finite subsets of $\Delta(K) \times \Delta(L)$. Assume all splittings on P and Q are available.

Example: $K = L = \{0, 1\}$, $P = Q = \{0, 1/3, 2/3, 1\}$ and

$$u =$$

$p = 1$	6	4	2	1
$p = 2/3$	4	2	4	3
$p = 1/3$	2	4	2	5
$p = 0$	0	2	4	7
	$q = 0$	$q = 1/3$	$q = 2/3$	$q = 1$

$$U = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & 2 & 4 & 3 \\ 2 & 4 & 2 & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}$$

$A = (A_{i,j})$ is S -concave if: $\forall j, \forall i \in \{2, 3\}, a_{i,j} \geq 1/2a_{i-1,j} + 1/2a_{i+1,j}$.

$A = (A_{i,j})$ is T -convex if: $\forall i, \forall j \in \{2, 3\}, a_{i,j} \leq 1/2a_{i,j-1} + 1/2a_{i,j+1}$.

$$\text{cav}U = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & 4 & 4 & 3 \\ 2 & 4 & 4 & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}, \quad \text{vex} \text{cav}U = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & 11/3 & 10/3 & 3 \\ 2 & 9/3 & 12/3 & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}$$

$$\text{vex}U = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & 2 & 5/2 & 3 \\ 2 & 2 & 2 & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}, \quad \text{cav} \text{vex}U = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & 10/3 & 8/3 & 3 \\ 2 & 8/3 & 10/3 & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}$$

How to compute $V = MZ(U)$, with $U = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & 2 & 4 & 3 \\ 2 & 4 & 2 & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}$?

First check first and last rows and columns: OK, $V = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & a & b & 3 \\ 2 & c & d & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}$

Guess $a > 2$, $b < 4$, $c < 4$ and $d > 2$ ($v > u$, $v < u$, $v = u$).

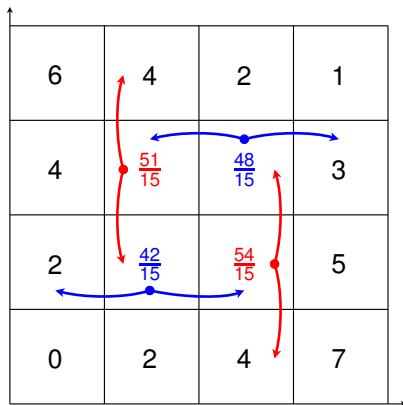
Solve $a = (4 + c)/2$, $b = (a + 3)/2$, $c = (2 + d)/2$ and $d = (b + 4)/2$.

Then

$$V = \begin{pmatrix} 6 & 4 & 2 & 1 \\ 4 & 51/15 & 48/15 & 3 \\ 2 & 42/15 & 54/15 & 5 \\ 0 & 2 & 4 & 7 \end{pmatrix}$$

→ finite algorithm

Optimal strategies:



The number of disclosure periods is unbounded, but disclosure stops with probability 1 in finite time (reminiscent of the 4 frogs in Forges, 1990).

(same phenomena can happen in the case of experiments.)

Approximation

Proposition 5: (Approximation) Consider the standard case with all splittings available and u continuous. Assume binary states.

For each $n \geq 1$, define U_n using uniform grids with stepsize $1/n$.

Compute V_n the MZ value of the matrix U_n .

Define $w_n : \Delta(K) \times \Delta(L)$ as the piecewise bilinear extension of V_n .

Then $(w_n)_n$ uniformly convergences to $MZ(u)$ when $n \rightarrow \infty$.

Splitting on a finite grid: general case

Assume P, Q finite and that non revealing splittings are available (no need for other assumption on S and T).

For all (p, q) in $A \times B$, define

$$v_-(p, q) = \sup\{v(s, q), \exists \alpha \in [0, 1), s \in \Delta(A \setminus \{p\}), \alpha \delta_p + (1 - \alpha)s \in S(p)\},$$

$$v_+(p, q) = \inf\{v(p, t), \exists \alpha \in [0, 1), t \in \Delta(B \setminus \{q\}), \alpha \delta_q + (1 - \alpha)t \in T(q)\}.$$

Proposition 6: (Long Information Design for general finite grids)

The games G_∞ and H_∞ have a value in pure strategies. This value is the unique function $v : A \times B \rightarrow \mathbb{R}$ which is S -concave, T -convex, and s.t. for all (p, q) in $A \times B$:

(Q1') if $v(p, q) > u(p, q)$, then $v(p, q) = v_-(p, q)$,

(Q2') if $v(p, q) < u(p, q)$, then $v(p, q) = v_+(p, q)$.

Remark: Pure ε -optimal strategies exist (use stopping times).

A few questions

Question 1: The conclusion of the theorem also hold by replacing A2 by A'2:

There exists a distance d on P such that: $\forall f$ 1-Lipschitz, $\text{cav}f$ is 1-Lipschitz.

Fix P a convex compact subset of $\Delta(K)$ and assume all splittings with posteriors in P are available. Can we find d such that A'2 holds ?

Question 2: All splittings available on $P = \Delta(K)$, but u is severely discontinuous: $u(p, q) = 1$ if $p - q$ is rational, $= -1$ otherwise.

Both $v = 1$ and $v = -1$ are S -concave, T -convex and satisfy (Q1) and (Q2).

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